

**Amendments to the Claims:**

This listing of claims will replace all prior versions, and listings, of claims in the application:

**Listing of Claims**

1.-7. (Cancelled)

8. (Currently Amended) An authentication process involving a first entitydevice, which possesses a public key v and a secret key s, the public and secret keys being related by an operation modulo n, where n is an integer, the modulus n being specific to the first entitydevice, and a second entitydevice, which knows the public key v, the first and second entities being provided with means to exchange zero-knowledge information and to carry out cryptographic calculations on the zero-knowledge information, calculations being carried out modulo n wherein in the process the modulo n operation is of  $v=s^{-t} \pmod{n}$ , t being a parameter and in that the modulo n calculations are performed according to the “Chinese remainders” method.

9. (Currently Amended) A process according to claim 8, wherein the information exchanges are of zero-knowledge and wherein the cryptographic calculations are completed as follows:

the first entitydevice selects are at least one integer r at random ranging between 1 and n-1 and calculates at least one parameter x equal to  $r^t \pmod{n}$ , then at least one number c that is at least one function of the at least one of a parameter and a message and sends the at least one number c to the second entitydevice;

the second entitydevice receives the at least number c, selects at least one number e at random, and sends the at least one number e to the first entitydevice;

the first entitydevice receives the at least one number e, carries out at least one calculation using the at least one number e and the secret key s, the result of the at least one calculation yielding at least one answer y, and sends the at least one answer y to the second entitydevice.

the second entitydevice receives the at least one answer y, carries out one calculation using the public key v and the modulus n, and checks with a modulo n operation that the result of the one calculation is coherent with the received at least one number c.

10. (Previously Presented) A process according to Claim 9, wherein a size of the number n, expressed in number of bits, is less than 1,000.

11. (Previously Presented) A process according to Claim 10, wherein a size of the number n is between 700 and 800.

12. (Cancelled)

13. (Currently Amended) A message signature process configured for a signatorydevice provided with a public key v and a secret key s, the public and private keys being related by a modulo n calculation, where n is an integer, which is specific to the signatorydevice, the process utilizing means configured to calculate at least one number c that is a function of a message M to be signed, configured to calculate at least one number y that is a function of the secret key s, and configured to transmit the numbers y and c that are the signature of the message and the message M, wherein the modulo n operation is  $v=s-t \pmod{n}$ , t being a parameter.

14. (Currently Amended) A message signature process according to claim 13, wherein the signatory device selects an integer  $r$  at random between 1 and  $n-1$ , calculates a parameter  $x$  equal to  $rt \pmod n$ , calculates at least one number  $e$  that is a function of parameter  $x$  and the message  $M$  to be signed, calculates the at least one number  $y$  using its secret key  $s$ , said at least one number  $y$  being a function of numbers  $r$  and  $e$ , and transmits the numbers  $c$  and  $y$  as the signature.